# Amalgamating many overlapping Boolean algebras

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#### Ternary obstructions to amalgamation

**Definition.** A sequence  $(A_i)_{i < n}$  of Boolean algebras is overlapping if, for all i, j, the Boolean operators of  $A_i$  and  $A_j$  agree when restricted to their common domain.

Given a pair of overlapping Boolean algebras A, B, there is a Boolean algebra C extending both of them.

Moreover, an arbitrary  $\Delta$ -system of overlapping Boolean algebras also has a common extension. (Koppelberg)

But three overlapping Boolean algebras (or just posets) A, B, C may not have a common extension. Minimal example:

$$x <_A y <_B < z <_C x$$
.

(Generate A from x, y and the relation  $x \land -y = 0$ ; similarly construct B and C.)

# Some direct limits need ternary amalgamation

A set D of sets is directed if each pair  $x, y \in D$  satisfies  $x \cup y \subset z$  for some  $z \in D$ .

**Proposition.** If D is a directed set of countable sets and  $|\bigcup D| \geq \aleph_n$ , then there are  $x_1, \ldots, x_n \in D$  such that  $\bigcap_{j \neq i} x_j \not\subset x_i$  for all i.

Therefore, any construction of a Boolean algebra of size  $\geq \aleph_3$  as a directed union of countable Boolean algebras must amalgamate non- $\Delta$ -system triples of overlapping algebras.

To ease such constructions, we combine:

- 1. Algebra: a sufficient condition for amalgamation.
- 2. Set theory: Long  $\omega_1$ -approximation sequences (also known as Davies sequences).

# Algebra: n-ary pushouts

**Definition.** A pushout of overlapping Boolean algebras  $(A_i)_{i < n}$  is a Boolean algebra  $\bigoplus_{i < n} A_i$  generated by:

- ▶ Distinct generators  $\bigoplus_i(x)$  for i < n and  $x \in A_i \setminus \{0_{A_i}, 1_{A_i}\}$ .
- Relations:
  - $\blacksquare_i(x \land y) = \boxplus_i(x) \land \boxplus_i(y) \text{ for } x, y \in A_i.$
  - $\blacksquare_i(-x) = \boxplus_i(x) \text{ for } x \in A_i.$
  - $\blacksquare_i(x) = \boxplus_j(x) \text{ if } x \in A_i \cap A_j.$

In the category of Boolean algebras and Boolean homomorphisms,  $\boxminus_{i < n} A_i$ , along with the morphisms  $\boxminus_i : A_i \to \boxminus_{i < n} A_i$ , is a colimit of the commutative diagram of inclusion maps id:  $\bigcap_{i \in s} A_i \to \bigcap_{i \in t} A_i$  for  $\varnothing \neq t \subset s \subset n$ .

# Algebra: *n*-wise commuting subalgebras

Notation:  $A \leq B$  means A is a subalgebra of B.

**Definition** (n = 2): Heindorf and Shapiro). Given  $A_i \leq B$  for i < n, we say  $(A_i)_{i < n}$  commutes in B if, for every tuple of ultrafilters  $U_i \in \text{Ult}(A_i)$  for i < n, if  $U_i \cap A_j = U_j \cap A_i$  for all i, j < n, then there is an ultrafilter  $V \in \text{Ult}(B)$  extending every  $U_i$ .

**Lemma.**  $(A_i)_{i < n}$  commutes in B iff we can choose  $\bigoplus_{i < n} A_i$  such that  $A_i \leq \bigoplus_{i < n} A_i \leq B$  for all i < n.

#### Application: An *n*-ary interpolation theorem

#### The Interpolation Theorem of Proposition Logic.

If  $\varphi \vdash \psi$ , then  $\varphi \vdash \chi \vdash \psi$  for some  $\chi$  with all its propositional variables common to  $\varphi$  and  $\psi$ .

The Interpolation Theorem can be reinterpreted as a corollary of certain pairs of subalgebras of a free Boolean algebra commuting.

#### An *n*-ary generalization.

If  $\bigwedge_{i \le n} \varphi_i \vdash \bot$ , then there exist  $\chi_i$  for i < n such that:

- $\triangleright \varphi_i \vdash \chi_i$  for each *i*.
- $ightharpoonup \bigwedge_{i < n} \chi_i \vdash \perp$ .
- For each i, each propositional variable in χ<sub>i</sub> is in φ<sub>i</sub> and in at least one other φ<sub>j</sub>.

# Algebra: a sufficient condition for amalgamation

Notation:  $\langle S \rangle$  denotes the Boolean closure of a subset S of a Boolean algebra.

**Theorem 1 (M., 2016).** Overlapping Boolean algebras  $(A_i)_{i < n}$  mutually extend to a pushout  $\bigoplus_{i < n} A_i$  if, for all  $k < m \le n$ ,

- 1.  $(A_i \cap A_m)_{i < m}$  commutes in  $A_m$ ,
- 2.  $(\coprod_i [A_i \cap A_m])_{i < m}$  commutes in  $\coprod_{i < m} A_i$ , and
- 3.  $\boxplus_k [A_k \cap A_m] = \boxplus_k [A_k] \cap \langle \bigcup_{i < m} \boxplus_i [A_i \cap A_m] \rangle$  in  $\bigoplus_{i < m} A_i$ .

It's not fun to verify all these conditions. Fortunately, there is a set-theoretic black box that hides these conditions behind one simpler condition.

#### Set theory: Long $\omega_1$ -approximation sequences

Let  $\mathfrak{H}$  be the structure  $(H(\theta), \in, \sqsubseteq_{\theta})$  where:

- lacktriangledown heta is a sufficiently large regular cardinal.
- ▶  $H(\theta)$  is the set of all sets hereditarily smaller than  $\theta$ .
- ightharpoonup  $\sqsubseteq_{\theta}$  well orders of  $H(\theta)$ .

#### Definition (M., 2008).

A transfinite sequence  $(M_{\alpha})_{\alpha<\eta}$  is a long  $\omega_1$ -approximation sequence if, for each  $\alpha$ :

- $M_{\alpha}$  is a countable elementary substructure of  $\mathfrak{H}$ .
- ▶ The sequence  $(M_{\beta})_{\beta < \alpha}$  is an element of  $M_{\alpha}$ .

**Lemma.** Given  $(M_{\alpha})_{\alpha<\eta}$  as above,

$$M_{\beta} \subsetneq M_{\alpha} \Leftrightarrow M_{\beta} \in M_{\alpha} \Leftrightarrow \beta \in \alpha \cap M_{\alpha}.$$

**Warning.**  $\{M_{\alpha} \mid \alpha < \eta\}$  is not a chain if  $\eta > \omega_1$ .

# Set theory: Coherence properties

**Lemma.** Given a long  $\omega_1$ -approximation sequence  $(M_\alpha)_{\alpha<\eta}$ :

For each  $\alpha < \eta$  and  $B \subset \eta$ , if  $M_{\alpha} \subset \bigcup_{\beta \in B} M_{\beta}$ , then  $M_{\alpha} \subset M_{\beta}$  for some  $\beta \in B$ .

For each nonempty  $S \subset \eta$ ,  $\bigcap_{\alpha \in S} M_{\alpha}$  is the *directed* union of of its subsets of the form  $M_{\beta}$ .

Each  $\alpha \leq \eta$  has a *finite* interval partition  $I_{\alpha}^{0}, \ldots, I_{\alpha}^{\neg(\alpha)-1}$  such that each  $\{M_{\beta} \mid \beta \in I_{\alpha}^{k}\}$  is *directed*.

If  $\alpha < \omega_n$ , then  $\Im(\alpha) \leq n$ ; if  $\alpha$  is a cardinal, then  $\Im(\alpha) = 1$ .

# Set theory: Pairing each $M_{\alpha}$ with a Boolean algebra

**Definition.** A Boolean  $\omega_1$ -complex is a sequence  $(A_\alpha, M_\alpha)_{\alpha < \eta}$  such that  $(M_\alpha)_{\alpha < \eta}$  is a long  $\omega_1$ -approximation sequence and, for all  $\alpha < \eta$ :

- 1.  $A_{\alpha}$  is a Boolean algebra.
- 2.  $A_{\alpha}$  is a subset of  $M_{\alpha}$ .
- 3.  $A_{\beta} \leq A_{\alpha}$  for all  $M_{\beta} \in M_{\alpha}$ .
- 4.  $A_{\alpha} \setminus \bigcup_{\beta < \alpha} A_{\beta}$  is disjoint from  $\bigcup_{\beta < \alpha} M_{\beta}$ .
- 5.  $(A_{\beta})_{\beta<\alpha}\in M_{\alpha}$ .
- 6.  $(A_{\alpha}^{k})_{k< \exists (\alpha)}$  commutes in  $A_{\alpha}$  where  $A_{\alpha}^{k} = \bigcup \{A_{\beta} \mid \beta \in I_{\alpha}^{k} \cap M_{\alpha}\}.$

Conditions 1–5 are trivial to satisfy provided  $\vec{A}$  and  $\vec{M}$  are constructed in parallel.

Condition 6 will guarantee that the sequence can be extended.  $\bigcup_{\alpha<\eta}A_\alpha \text{ is a directed union if } \eta \text{ is a cardinal.}$ 

#### Set theory: an easier amalgamation theorem

#### Theorem 2 (M., 2016.) If:

- $(A_{\alpha}, M_{\alpha})_{\alpha < \eta}$  is a Boolean  $\omega_1$ -complex,
- $(M_{\alpha})_{\alpha<\eta+1}$  is a long  $\omega_1$ -approximation sequence, and
- $\blacktriangleright (A_{\alpha})_{\alpha<\eta}\in M_{\eta},$

then  $B=\coprod_{k<\daleth(\eta)}A_{\eta}^{k}$  extends  $A_{\alpha}$  for all  $M_{\alpha}\in M_{\eta}.$ 

Therefore, to extend to a longer Boolean  $\omega_1$ -complex  $(A_\alpha,M_\alpha)_{\alpha<\eta+1}$ , we may choose any  $A_\eta$  meeting the following requirements.

- ▶  $B \leq A_{\eta}$ .
- $A_{\eta}$  is a subset of  $M_{\eta}$ .
- ▶  $A_{\eta} \setminus \bigcup_{\alpha < \eta} M_{\alpha}$  is disjoint from  $\bigcup_{\alpha < \eta} M_{\alpha}$ .

# Application: a higher-arity Freese-Nation property

#### Definition.

- ▶ Given  $B \le A$ , we say B is relatively complete in A and write  $B \le_{\rm rc} A$  if for every  $x \in A$  the set  $\{y \in B \mid y \le x\}$  has a maximum element.
- ▶ A Boolean algebra A has the n-ary FN if there is a club  $\mathcal C$  of countable subalgebras of A such that  $\langle B_1 \cup \cdots \cup B_{n-1} \rangle \leq_{\mathrm{rc}} A$  for all  $B_1, \ldots, B_{n-1} \in \mathcal C$ .
- ▶ A Boolean algebra A is projective if it is a retract of some free Boolean algebra F. (Retract means  $A \leftarrow_r F \leftarrow_e A$ ;  $r \circ e = id$ )

#### Theorem 3 (M., 2016).

- ▶ A is projective iff it has the *n*-ary FN for all *n*.
- ▶ If  $|A| < \aleph_n$  and A has the n-ary FN, then A is projective.
- ▶ For each n, there is a Boolean algebra of size  $\aleph_n$  with the n-ary FN but without the (n+1)-ary FN.

# Application: a higher-arity strong Freese-Nation property

**Definition** (n = 2: Heindorf and Shapiro).

A Boolean A has the n-ary strong FN if it has a cofinal family  $\mathcal C$  of *finite* subalgebras such that  $B_1, \ldots, B_n$  commutes in A for  $B_1, \ldots, B_n \in \mathcal C$ .

#### Theorem 4 (M., 2016).

- ▶ The *n*-ary strong FN implies the *n*-ary FN.
- ▶ A is projective iff it has the *n*-ary strong FN for all *n*.
- ▶ If  $|A| < \aleph_n$  and A has the n-ary strong FN, then A is projective.

#### Finitary applications

Let F be Stone dual of the Vietoris hyperspace functor or a nontrivial symmetric power functor.

F destroys the projectivity of the free Boolean algebra of size  $\aleph_2$ . (Ščepin)

**Corollary (M., 2016).** There is a *finite* Boolean algebra A with subalgebras  $B_1$ ,  $B_2$ ,  $B_3$  that commute in A but  $F(B_1)$ ,  $F(B_2)$ ,  $F(B_3)$  do not commute in F(A).

The above corollary is non-constructive and gives no bound on the size of A. One of my students, René Montemayor, found that the minimal A is  $\mathcal{P}(4)$ .

#### Open problems

- To what extent do the amalgamation theorems 1 and 2 generalize to arbitrary categories? At minimum, we must assume the category has limit and colimits of all finite diagrams.
- For all  $n \ge 1$ , the n-ary FN does not imply the (n + 1)-ary FN. For the strong FN, this is only known for n = 1, 2. Is the 4-ary strong FN stricter stronger than the 3-ary strong FN?
- The binary strong FN is known to be strictly stronger than the binary FN. (M., 2014) Is the ternary strong FN strictly stronger than the ternary FN?
- What is the algorithmic complexity of deciding a given list of overlapping finite Boolean algebras, reasonably encoded in N bits, has a common extension? A brute force search algorithm gives upper bounds of  $\mathrm{CoNP^{NP}}$  and space complexity  $O(\sqrt{N})$ .

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